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Transformational semantics for Evolving Logic Programs

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March 30, 2007

1 Introduction

Logic programming Evolving logic programs

2 Transformational semantics for EVOLP

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1 Introduction

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Logic program

 $\begin{array}{l} day \leftarrow \text{not } night.\\ night \leftarrow \text{not } day.\\ play_violin \leftarrow day.\\ play_piano \leftarrow night. \end{array}$



Logic program

 $day \leftarrow \text{not } night.$ $night \leftarrow \text{not } day.$ $play_violin \leftarrow day.$ $play_piano \leftarrow night.$

Stable models

 $M_1 = \{ day, play_violin \}$ $M_2 = \{ night, play_piano \}$

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- language based on logic programming
- intended for dynamic environments (e.g. multiagent systems)

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• syntax is the same as the syntax of logic programs

$$P = \begin{cases} write_thesis \leftarrow \text{ not } tired.\\ drink_coffee \leftarrow tired, \text{ not } no_coffee.\\ buy_coffee \leftarrow tired, no_coffee.\\ assert(tired) \leftarrow write_thesis.\\ assert(\text{not } tired) \leftarrow drink_coffee. \end{cases}$$

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Time	Program	Event	Model
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P











$$E_1 \quad P \longrightarrow M_1$$







$$P_2 = \{r \mid assert(r) \in M_1\}$$





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Evolution stable model $P_2 = \{r \mid assert(r) \in M_1\}$ $P_3 = \{r \mid assert(r) \in M_2\}$ $P_n = \{r \mid assert(r) \in M_{n-1}\}$

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Evolving logic program



Evolution stable models

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• shows that EVOLP can be, in fact, translated into traditional logic programming (and how)

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- shows its semantics from a different perspective
- provides a simple, straight-foward way to implement it

- definition of the transformation
- proofs of its correctness and completeness
- an implementation of propositional EVOLP

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How big is the transformed program?

For an input program *P* and an event sequence $\mathcal{E} = (E_1, E_2, \dots, E_n)$ over the universe \mathcal{U} we have:

- transformed program is built of $2n|\mathcal{U}|$ atoms (at most)
- lower bound for the size of the transformed program:

$$|P_{\mathcal{E}}| \ge n|P| + \sum_{k=1}^{n} |E_k|$$

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• upper bound for the size of the transformed program:

$$|P_{\mathcal{E}}| \le \frac{7}{2} \left(|P| \frac{n^3 + 5n}{6} + \sum_{k=1}^n |E_k| \left(\frac{(n-k)^3 + 5(n-k)}{6} + 1 \right) \right) + n|\mathcal{U}|$$

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• upper bound if input programs contain no nested asserts:

$$|P_{\mathcal{E}}| \leq \frac{7}{2} \left(|P| \frac{n^2 + n}{2} + \sum_{k=1}^n |E_k| (n - k + 1) \right) + n |\mathcal{U}|$$

- extensions of the existing implementation:
 - variable support
 - support for arithmetic predicates
 - ...
- a different, more direct implementation that could be used on-line

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Thank you. Are there any questions?

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