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**Introduction  
of dyadic  
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The Schema  
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# Introduction of dyadic pairing into PA

## Lecture 9

# Review: Dyadic concatenation

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$$x \mid y \leftrightarrow \exists z x \cdot z = y$$

$$Pow_2(p) \leftrightarrow \forall d ((d + 2) \mid p \rightarrow 2 \mid d)$$

$$2^{|x|} = \mu_p [Pow_2(p) \wedge x + 1 < 2 \cdot p]$$

$$x \star y = x \cdot 2^{|y|} + y$$

We have also defined

$$D\_two(m) \leftrightarrow \exists n_1 \exists n_2 m = n_1 \star 2 \star n_2 ,$$

i.e.  $\neg D\_two(m)$  iff  $m = 1 \star \dots \star 1$ .

However, we have for  $x > 0$   $2^{|x|} = \overbrace{1 \star \dots \star 1 \star 2}^{|x|}$

Also  $2^{|0|} = 1$ . Thus, for any  $x$   $2^{|x|} \dot{-} 1 = \overbrace{1 \star \dots \star 1}^{|x|}$  i.e.

$$\neg D\_two(m) \leftrightarrow \exists x m + 1 = 2^{|x|} \leftrightarrow \exists x m = 2^{|x|} \dot{-} 1$$

# Marker sequences

To every number  $m$  there are unique numbers  $n$  and  $b$  as well as numbers  $a_1, a_2, \dots, a_n$  of unique dyadic length such that

$$m = \overbrace{1 \star \dots \star 1 \star 2 \star 1 \star \dots \star 1 \star 2 \star \dots \star 1 \star \dots \star 1 \star 2}^n \star \overbrace{1 \star \dots \star 1}^b$$

This can be also written as:

$$m = 2^{|a_1+1|} \star 2^{|a_2+1|} \star \dots \star 2^{|a_n+1|} \star (2^{|b|} - 1)$$

Note that that  $x + 1$  and  $2^{|x+1|}$  have the same **length** because of **idempotency**:  $2^{|x+1|} = 2^{|2^{|x+1|}|}$ .

To every  $m$  and  $w$  of the same **length** there are **unique**  $n, b, a_1, a_2, \dots, a_n$  s.t.:

$$w = (a_1 + 1) \star (a_2 + 1) \star \dots \star (a_n + 1) \star b$$

$$m = 2^{|a_1+1|} \star 2^{|a_2+1|} \star \dots \star 2^{|a_n+1|} \star (2^{|b|} - 1)$$

Square is an **injection**:

$$x^2 = y^2 \rightarrow x = y$$

and  $\sqrt{2}$  is irrational:  $y > 0 \rightarrow (\frac{x}{y})^2 \neq 2$  for  $x, y \in \mathbb{Z}$ . This is expressed in PA as

$$x^2 = 2 \cdot y^2 \rightarrow x = 0 \wedge y = 0$$

Hence

$$i \cdot x^2 = j \cdot y^2 \wedge x > 0 \wedge 0 < i, j \leq 2 \rightarrow i = j \wedge x = y$$

This property is used in the **uniqueness** property on the following slide.

## Existence of splits

$$\exists w \exists i \exists m (t = w \star i \star m \wedge 2^{|w|} = 2^{|m|} \wedge 2^{|i|} \leq 2)$$

We can view this as every number  $t$  can be decomposed into

two numbers of length almost  $\sqrt{|t|}$ :  $t = \begin{array}{|c|c|} \hline w & i \\ \hline \end{array} \star \begin{array}{|c|c|} \hline m & \\ \hline \end{array}$  and  $i = 0, 1, 2$

## Uniqueness of splits:

$$w_1 \star i_1 \star m_1 = w_2 \star i_2 \star m_2 \wedge 2^{|i_1|} \leq 2 \wedge 2^{|w_1|} = 2^{|m_1|} \wedge$$

$$2^{|i_2|} \leq 2 \wedge 2^{|w_2|} = 2^{|m_2|} \rightarrow w_1 = w_2 \wedge i_1 = i_2 \wedge m_1 = m_2$$

# Definition of splits

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$$t \doteq \begin{bmatrix} v \\ m \end{bmatrix} \leftrightarrow t = v \star m \wedge \exists w \exists i (v = w \star i \wedge 2^{|m|} = 2^{|w|} \wedge 2^{|i|} \leq 2)$$

We can visualize the definition as:  $t = \begin{array}{|c|c|} \hline w & i \\ \hline v \\ \hline \end{array} \star \begin{array}{|c|c|} \hline m & \\ \hline \end{array}$ .

Splits **exist**  $\exists v \exists m t \doteq \begin{bmatrix} v \\ m \end{bmatrix}$  and they are **unique**:

$$t \doteq \begin{bmatrix} v_1 \\ m_1 \end{bmatrix} \wedge t \doteq \begin{bmatrix} v_2 \\ m_2 \end{bmatrix} \rightarrow v_1 = v_2 \wedge m_1 = m_2$$

# Adjustment of splits

If  $t \doteq \begin{bmatrix} v \\ m \end{bmatrix}$  then  $t = w \star i$  where  $i = 0, 1, 2$  and

$$w = (a_1 + 1) \star (a_2 + 1) \star \dots \star (a_n + 1) \star b$$

$$m = 2^{|a_1+1|} \star 2^{|a_2+1|} \star \dots \star 2^{|a_n+1|} \star (2^{|b|} - 1)$$

The **tail** part such that:  $b \star i \star (2^{|b|} - 1) \doteq \begin{bmatrix} b \star i \\ 2^{|b|} - 1 \end{bmatrix}$  is an **atom**.

We now define  $Adj(t) \doteq \begin{bmatrix} w' \\ m' \end{bmatrix}$  removing the atom:

$$w' = (a_1 + 1) \star (a_2 + 1) \star \dots \star (a_n + 1)$$

$$m' = 2^{|a_1+1|} \star 2^{|a_2+1|} \star \dots \star 2^{|a_n+1|}$$

$$Adj(t) = s \leftrightarrow \exists v_1 \exists v_2 \exists m \exists b (t \doteq \begin{bmatrix} v_1 \star v_2 \\ m \mathbf{0} \star (2^{|b|} - 1) \end{bmatrix} \wedge 2^{|v_1|} = 2^{|m \mathbf{0}|} \wedge$$

$$s \doteq \begin{bmatrix} v_1 \\ m \mathbf{0} \end{bmatrix})$$

Note that:  $v_1 = w'$ ,  $m \mathbf{0} = m'$ , and  $v_2 = b \star i$

We define the **concatenation**:

$$s \boxplus t = u \leftrightarrow \exists v_1 \exists v_2 \exists m \exists x (Adj(s) \dot{=} \begin{bmatrix} v_1 \\ m_1 \end{bmatrix} \wedge t \dot{=} \begin{bmatrix} v_2 \\ m_2 \end{bmatrix} \wedge \\ u \dot{=} \begin{bmatrix} v_1 \star v_2 \\ m_1 \star m_2 \end{bmatrix})$$

and the **pairing**:

$$x; t = ((x + 1) \star 2^{|x+1|}) \boxplus t$$

Thus

$$t \dot{=} \begin{bmatrix} v \\ m \end{bmatrix} \rightarrow s; t \dot{=} \begin{bmatrix} (x + 1) \star v \\ 2^{|x+1|} \star m \end{bmatrix}$$