

Symbolic programming with dyadic concatenation

Dyadic concatenation

Recall that for every $x \in \mathbb{N}$ **exactly one** holds:

$$0 \quad x = y\mathbf{1} = 2 \cdot y + 1 \quad x = y\mathbf{2} = 2 \cdot y + 2$$

Two **strings of dyadic digits**, say 1211 and 2122 can be **concatenated** to

12112122 = 1211 \star 2122 with the function:

$$x \star 0 = y \quad x \star y\mathbf{1} = (x \star y)\mathbf{1} \quad x \star y\mathbf{2} = (x \star y)\mathbf{2}$$

Dyadic concatenation is associative and

$$x \star y = x \cdot 2^{|y|_d} + y$$

where the **dyadic size** function $|x|_d$ is:

$$|0|_d = 0 \quad |x\mathbf{1}|_d = |x\mathbf{2}|_d = |x|_d + 1$$

We can also define $Drop(n, x)$ and $Take(n, x)$ s.t.:

$$n \leq |x|_d \rightarrow x = Drop(n, x) \star Take(n, x)$$

Dyadic Coding of Turing Machines

Turing has invented his **machines** to investigate **computability**. Such a machine M can be seen as working on a **tape** consisting of symbols 1 and 2. M starts with the tape containing **input** and it terminates with the tape containing **output**. For instance:

$$11112111\underline{2} \xrightarrow{M} 1111111\underline{2}$$

M can be viewed as performing **addition** $1^4 + 1^3 = 1^7$ in **monadic** where the symbol 2 separates the numbers. And $\underline{\quad}$ designates the **current** position of its **reading-writing head**.

The machine executes a **program** with **instructions**: l (move the head one position **left**), **extending** the tape with the symbol 2 if the head is at the left end,
 r (move one position **right**), extending with 2 if needed,
 w_1 (**overwrites** the currently read symbol with 1);
 w_2 (**overwrites** with 2),
 $If_1(p_1, p_2)$ executes the **program** p_1 if the currently **read** symbol is 1 and the **program** p_2 if the symbol is 2
 $Wh_1(p)$ executes the program p $Wh_1(p)$ if the currently **read** symbol is 1 and does **nothing** if the symbol is 2 .

Program $Rb \equiv r Wh_1(r)$ moves **one block right**:

$$\dots 1^m \underline{2} 1^n \underline{2} \dots \mapsto \dots 1^m 2 1^n \underline{2} \dots$$

Program $Lb \equiv l Wh_1(l)$ moves **one block left**:

$$\dots 1^m 2 1^n \underline{2} \dots \mapsto \dots 1^m \underline{2} 1^n 2 \dots$$